## Extended Haagerup exists!

#### Scott Morrison http://tqft.net/ joint work with Stephen Bigelow, Emily Peters and Noah Snyder

Microsoft Station Q / UC Santa Barbara

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# Outline

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  - Subfactors and ⊗-categories
  - Planar algebras: pictures for subfactors

2 Haagerup's classification up to index  $3 + \sqrt{3}$ 

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  - Existence
  - Construction
- 4 Skein theory for extended Haagerup
  - Quadratic relations
  - The jellyfish algorithm

 $\begin{array}{l} \mbox{Subfactors for native speakers of $\otimes$-categories$}\\ \mbox{Haagerup's classification up to index 3 + $\sqrt{3}$}\\ \mbox{Constructing planar algebras}\\ \mbox{Skein theory for extended Haagerup} \end{array}$ 

Subfactors and  $\otimes$ -categories Planar algebras: pictures for subfactors

## Subfactors and ⊗-categories

- Extremal finite-index *II*<sub>1</sub> subfactors correspond to unitary spherical ⊗<sub>A</sub>, ⊗<sub>B</sub>-categories with a chosen generator <sub>A</sub>X<sub>B</sub>.
  - (objects are bimodules for A ⊂ B, and X is the 'regular' bimodule <sub>A</sub>B<sub>B</sub>)
- The 'index' [A:B] of the subfactor is  $(\dim_q X)^2$ .
- The 'principal graph' encodes the tensor products V ⊗<sub>A</sub> X and V ⊗<sub>B</sub> X\*.
- The 'even part' is the subcategory of *A*-*A* objects (alternatively of *B*-*B* objects).
- The double of the even part is a modular tensor category.
  - (it doesn't matter which even part you take, because they're Morita equivalent)
  - (these MTCs may be 'exotic', that is, don't come from quantum groups [Hong, Rowell, Wang])

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Subfactors and  $\otimes$ -categories Planar algebras: pictures for subfactors

# Planar algebras: pictures for subfactors



- To produce the planar algebra, restrict your objects to  $X \otimes X^* \otimes X \otimes \cdots$ .
- To recover the  $\otimes_A$ ,  $\otimes_B$ -category, form the idempotent completion.

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# Haagerup's list

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# Haagerup's list

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- Haagerup and Asaeda & Haagerup (1999) constructed two of these possibilities.
- Bisch (1998) and Asaeda & Yasuda (2007) ruled out infinite families.
- Today, we construct the missing example ('extended Haagerup'), and complete the classification. • = • • = •

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- Asaeda & Haagerup's constructions find a biunitary flat connection, using Ocneanu's paragroup formalism.
- Ikeda (1998) did numerical calculations suggesting that the extended Haagerup subfactor existed "up to  $10^{-9}$ ".
- Our techniques are very different. Although we made heavy use of a computer while searching for a solution, verifying the answer only requires a computer to multiply large matrices, using exact arithmetic.

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Annular and quadratic tangles Existence Construction

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Annular and quadratic tangles Existence Construction

#### Project

Now that we know about planar algebras, we can give <u>skein theory</u> constructions of subfactors.

- Jones' work on <u>annular</u> and <u>quadratric</u> tangles gives you information about the generators and relations to expect.
- Peters [math/0902.1294] has implemented this approach to give a new construction of the Haagerup subfactor. Our work follows that model closely.

Annular and quadratic tangles Existence Construction

#### Annular Temperley-Lieb

Every planar algebra  $\mathcal{P}$  is a representation of the <u>annular</u> <u>Temperley-Lieb</u> category. Every representation breaks up as a sum of representations generated by <u>lowest weight vectors</u>  $T_{k,\lambda} \in \mathcal{P}_k$ satisfying



for some  $k \ge 0$  and  $\lambda \in \sqrt[k]{1}$ . Thus the planar algebra is generated by elements of this form.

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Annular and quadratic tangles Existence Construction

If T is a lowest weight vector in  $\mathcal{P}_5$ , then we write  $\mathcal{ATL}_{+1}(T)$  for the "annular consequences" of T in  $\mathcal{P}_6$ .



#### Lemma

If  $\delta > 2$ , the annular consequences  $\mathcal{ATL}_{+k}(T)$  of a lowest weight vector T are all linearly independent.

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Annular and quadratic tangles Existence Construction

## Quadratic tangles

Further, by counting the dimensions of the  $\mathcal{ATL}$  representations, we can expect relations between certain <u>quadratic tangles</u>.

#### Example (Extended Haagerup)

$$\mathcal{P}\cong V_{0,\delta}\oplus V_{8,-1}\oplus V_{10,\bullet}\oplus\cdots$$

If T generates  $V_{8,-1}$ , then

- $\underline{}^{8}(T) \underline{}^{8}(T) \underline{}^{8} \in P_{8}$  is a linear combination of T and  $\mathcal{TL}$ ,
- there no lowest weight vector in  $P_9$ , so  $\xrightarrow{9}{}$   $T \xrightarrow{7}{}$   $T \xrightarrow{9}{}$  must be a linear combination of  $\mathcal{ATL}_{+1}(T)$  and  $\mathcal{TL}$ ,
- there's only one lowest weight vector in  $P_{10}$ , so some linear combination of  $\frac{10}{T} + \frac{10}{T} + \frac{10}{10}$  and  $\frac{10}{T} + \frac{10}{10} + \frac{10}{10}$  must lie in  $\mathcal{ATL}_{+2}(T) \oplus \mathcal{TL}$ .

Annular and quadratic tangles Existence Construction

## Existence

#### Question

Consider a planar algebra  $\mathcal{P}$  generated by an element T satisfying relations like these. Is it the extended Haagerup planar algebra?

- In fact, if we can show that P is a <u>subfactor planar algebra</u> with the correct index (the largest root of x<sup>3</sup> − 8x<sup>2</sup> + 17x − 5, ~ 4.3772...), then Haagerup's classification guarantees it must have the desired principal graph.
- It's probably also possible to compute the principal graph directly from the skein theory.

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Annular and quadratic tangles Existence Construction

## When is $\mathcal{P}$ a subfactor planar algebra?

- Do we know the relations are consistent? (i.e., is  $\mathcal{P} \neq 0$ ?)
- Is the planar algebra unitary?
- Is dim  $\mathcal{P}_0=1?\,$  That is, can we evaluate every closed diagram using the relations?

To answer the first two questions, we can try to find an element T inside a larger unitary planar algebra. Fortunately, there's an obvious place to look:

#### Theorem

Every subfactor planar algebra  $\mathcal{P}$  is a subalgebra of the graph planar algebra of the principal graph  $\Gamma(\mathcal{P})$ .

The evaluation problem requires some clever skein theory.

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Annular and quadratic tangles Existence Construction

## The graph planar algebra

#### Definition

The graph planar algebra  $GPA(\Gamma)$  of a bipartite graph  $\Gamma$  has spaces  $GPA(\Gamma)_k = \mathbb{C} \{ \text{length } 2k \text{ loops on } \Gamma \}$  and  $\mathcal{TL}$  action described in Jones', "The planar algebra of a bipartite graph", in terms of the Perron-Frobenius eigenvector.

The graph planar algebra is always  $\underline{unitary}$  and  $\underline{spherical}$ . It is not a subfactor planar algebra since

 $\dim(\mathit{GPA}(\Gamma)_0) = \sharp \text{even vertices} > 1.$ 

# Example $\dim(GPA(\Gamma(\mathcal{H}))_4) = 375, \dim(GPA(\Gamma(\mathcal{EH}))_8) = 148475$ Scott Morrison Extended Haagerup exists!

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## Construction

We now start solving equations in  $GPA(\Gamma(\mathcal{EH}))_{16}$ .

- Considerations from Jones' Quadratic Tangles tells us the rotational eigenvalue of the generator must be -1.
- We cut down  $GPA(\Gamma(\mathcal{EH}))_{16}$  to a 19 dimensional space using the linear relations



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- The generator must satisfy  $T^2 = f^{(16)}$  this gives 148475 quadratic equations!
- We look at specially chosen subset of these and find a solution T by ad-hoc methods.

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At this point you can forget everything I told you about finding T!

- We can write down an explicit description of *T* as an element of the 148475*d* vector space. See the generator!
- It's tedious but straightforward to check by hand that T is a lowest weight vector with rotational eigenvalue -1.
- Certainly the subalgebra generated by *T* is both nontrivial and unitary.
- We still need to see that all closed diagrams can be evaluated, that is,  $\mathcal{P}(T)_0 = \mathbb{C}$ .

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Quadratic relations The jellyfish algorithm

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Subfactors for native speakers of  $\otimes$ -categories Constructing planar algebras Skein theory for extended Haagerup

Quadratic relations

# Quadratic relations

with

We have a candidate generator; let's find the particular quadratic relations it satisfies. First, with the aid of a computer (doing exact arithmetic!) we calculate six moments.

tr 
$$(T^2) = [9] \sim 24.66097...$$
  
tr  $(T^3) = 0$   
tr  $(T^4) = [9]$   
tr  $(\rho^{1/2}(T)^2) = -[9]$   
tr  $(\rho^{1/2}(T)^3) = -i\frac{[18]}{\sqrt{[8][10]}} \sim -15.29004i$   
tr  $(\rho^{1/2}(T)^4) = \frac{1}{5} (46\lambda^4 - 2\lambda^2 - 94) \sim 34.1409...$   
with  $\lambda$  the root of  $x^6 + 2x^4 - 3x^2 - 5$  which is approximately  
1.12867i

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Bessel's inequality now gives the desired relations



Note that the first equation only holds with the given shading. For the Haagerup and extended Haagerup graph, when we write  $\frac{9}{T}T^{T}T^{9}$  in terms of  $\mathcal{ATL}_{+1}(T)$  a certain coefficient is 'unexpectedly' zero.

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Now substitute the first equation into the last term of the second, and expand the Jones-Wenzl idempotents.



This relation lets us 'pull a generator through a pair of strands'. This increases the number of generators in the diagram, so it isn't immediately obvious how this helps evaluate closed diagrams!

Quadratic relations The jellyfish algorithm

#### Theorem

Stephen's jellyfish algorithm shows these relations suffice to evaluate arbitrary closed diagrams.



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#### Begin with arbitrary planar network of Ts.



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Now float each generator to the surface, using the relation. =

Quadratic relations The jellyfish algorithm

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Quadratic relations The jellyfish algorithm

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Quadratic relations The jellyfish algorithm

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Now float each generator to the surface, using the relation.

Quadratic relations The jellyfish algorithm

The diagram now looks like a polygon with some diagonals, labelled by the numbers of strands connecting generators.



- Each such polygon has a corner, and the generator there is connected to one of its neighbours by at least 8 edges.
- Use  $T^2 = f^{(16)}$  to reduce the number of generators, and recursively evaluate the entire diagram.

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Quadratic relations The jellyfish algorithm

## Thank you!



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An explicit generator

To specify the generator T we need to give its value on every one of the 148475 loops of length 16.

$$T(\gamma) = r \cdot \sigma(\gamma) \cdot p_{\widehat{\gamma}} \cdot rac{1}{d_{\gamma_1}} \cdot \prod_{i=1}^{16} rac{1}{\sqrt{d_{\gamma_i}}}.$$

• 
$$r = -1843700 + 5847375d^2 - 1614050d^4$$

- If  $\gamma$  is a path on the principal graph, produce a sequence  $\widehat{\gamma} \in \{0, 1, 2\}^8$  so that if  $\gamma_{2i-1}$  is in the *j*-th arm of the principal graph, then  $\widehat{\gamma}_i = j$ .
- Let  $\sigma(\gamma)$  to be -1 raised to the number of times the vertices  $v_0, w_0, z_i$  and  $a_i$  appear in  $\gamma$ .

We still need to specify  $3^8 = 6561$  values for  $p_{\widehat{\gamma}}$ .

An explicit generator

Fix  $\lambda = \sqrt{(2-d^2)}$ . Define 21 elements of  $\mathbb{Z}[\lambda]$ 

 $p_{00000001} = -2\lambda^4 - \lambda^2 + 9$  $p_{00000101} = 2\lambda^4 + \lambda^2 = 0$  $p_{00001001} = \lambda^5 - \lambda^3 - 3\lambda$  $p_{00010001} = 2\lambda^4 + \lambda^2 - 9$  $p_{00010101} = \lambda^4 - 2\lambda^2 + 1$  $p_{00011011} = \lambda^4 - \lambda^2 - 3$  $p_{00100111} = \lambda^2 + 1$  $p_{00101101} = \lambda^5 - \lambda$  $p_{00110111} = -\lambda^5 - 2\lambda^3 - 4\lambda^2 - \lambda - 5$  $p_{01010111} = \lambda^4 + \lambda^2$  $p_{01110111} = \lambda^4 + 6\lambda^2 + 6$ 

 $p_{00000011} = -\lambda^5 - \lambda^3 + 3\lambda$  $p_{00000111} = 1$  $p_{00001011} = \lambda^3 - 1$  $p_{00010011} = \lambda^5 - \lambda^4 + \lambda^2 - 3\lambda + \lambda^2$  $p_{00010111} = 1 - \lambda^4$  $p_{00100101} = 5 - 2\lambda^4$  $p_{00101011} = -\lambda^5 - \lambda^3 + \lambda + 1$  $p_{00110011} = 2\lambda^5 + 5\lambda^3 + 4\lambda$  $p_{01010101} = -4\lambda^4 + 3\lambda^2 + 7$  $p_{01011011} = \lambda^4 - 2\lambda^2 - 4$ 

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#### Extend these definitions to every $p_w$ for $w \in \{0,1\}^8$ by the rules

 $p_{abcdefgh} = -p_{bcdefgha}$  $p_{abcdefgh} = \overline{p}_{ahgfedcb}$ 

and

 $p_{00000000} = 0$  $p_{abcd1111} = 0.$ 

Note that we have to check these rules are well-defined. For example, one can get from  $p_{00110011}$  to  $p_{01100110}$  either by rotating, or by reversing.

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# Further extend these definitions to every $p_w$ for $w \in \{0, 1, 2\}^8$ by the rules

$$p_{x0y} + p_{x1y} + p_{x2y} = 0.$$

◄ Return to the talk...

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